

Why concept lattices are large

Extremal theory for the number of minimal generators and formal concepts

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Abstract. A unique type of subcontexts is always present in formal contexts with many concepts: the contranominal scales. We make this precise by giving an upper bound for the number of minimal generators (and thereby for the number of concepts) of contexts without contranominal scales larger than a given size. Extremal contexts are constructed which meet this bound exactly. They are completely classified.

1 Introduction

The primitive data model of Formal Concept Analysis is that of a formal context, which is unfolded into a concept lattice for further analysis. It is well known that concept lattices may be exponentially larger than the contexts which gave rise to them. An obvious example is the boolean lattice $B(k)$, having 2^k elements, the standard context of which is the $k \times k$ *contranominal scale* $\mathbb{N}^c(k)$. This is not the only example of contexts having large associated concept lattices: indeed, the lattice of any subcontext is embeddable in the lattice of the whole context [4], which means that contexts having large contranominal scales as subcontexts necessarily have large concept lattices as well. Those considerations induce one natural question, namely, whether there are other reasons for a concept lattice to be large. As it will be shown in this paper, the answer is no.

The structure of the paper is as follows. Our starting point is a known upper bound for the number of concepts, which we improve using the language of minimal generators. Then, we show that our result is the best possible by constructing lattices which attain exactly the improved upper bound. These lattices, i.e., the extremal lattices, are characterized.

2 Fundamentals

For a set S , a context of the form (S, S, \neq) will be called a *contranominal scale*. We will denote by $\mathbb{N}^c(k)$ the contranominal scale with k objects (and k attributes), that is, the context $([k], [k], \neq)$, where $[k] := \{1, 2, \dots, k\}$. The expression $\mathbb{K}_1 \leq \mathbb{K}$ denotes that \mathbb{K}_1 is a subcontext of \mathbb{K} . The symbol \cong expresses the existence of an order-isomorphism whenever two ordered sets are involved

or, alternatively, the existence of a context isomorphism in the case of formal contexts. For a context \mathbb{K} to be $\mathbb{N}^c(k)$ -free means that there does not exist a subcontext $\mathbb{K}_1 \leq \mathbb{K}$ with $\mathbb{K}_1 \cong \mathbb{N}^c(k)$. The boolean lattice with k atoms, that is, $\mathfrak{B}(\mathbb{N}^c(k))$, will be denoted by $B(k)$. Similarly, we say that a lattice L is $B(k)$ -free whenever $B(k)$ does not (order-)embed into L . Using Proposition 32 from [4] one has that \mathbb{K} is $\mathbb{N}^c(k)$ -free whenever $\mathfrak{B}(\mathbb{K})$ is $B(k)$ -free. The converse is also true and is, in fact, the content of our first proposition. An example of a context which has $\mathbb{N}^c(3)$ as a subcontext along with its concept lattice is depicted in Figure 1. One may observe that the context is $\mathbb{N}^c(4)$ -free, since its lattice has ten concepts (and would have at least sixteen otherwise).

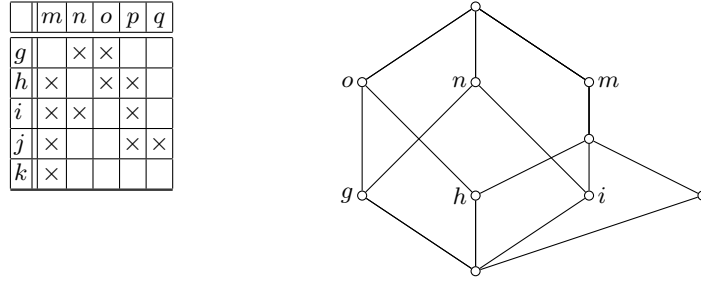


Fig. 1: A context \mathbb{K} with $\mathbb{N}^c(3) \leq \mathbb{K}$ and its concept lattice. The object and attribute concepts belonging to the $B(3)$ suborder are indicated on the diagram.

We denote by $J(L)$ and $M(L)$, respectively, the set of completely join-irreducible and meet-irreducible elements of a lattice L . The least and greatest elements of a lattice L will be denoted, respectively, by 0_L and 1_L . The symbol \prec will designate the covering relation between elements, that is, $x \prec y$ if $x < y$ and for every $z \in L$, $x < z \leq y \Rightarrow z = y$. The *length* of a finite lattice is the number of elements in a maximum chain minus one. An *atom* is an element covering 0_L , while a *coatom* is an element covered by 1_L . Whenever two elements $x, y \in L$ are incomparable, we will write $x \parallel y$. We denote by $A(L)$ the set of atoms of a lattice L . For an element $l \in L$, we shall write $\downarrow l := \{x \in L \mid x \leq l\}$ as well as $\uparrow l := \{x \in L \mid x \geq l\}$. Moreover, for $l \in L$ we denote by A_l the set $A(L) \cap \downarrow l$ and, similarly, by J_l the set $J(L) \cap \downarrow l$. A complete lattice L is called *atomistic* if $x = \bigvee A_x$ holds for every $x \in L$. In this case, $A(L) = J(L)$.

Proposition 1. *Let \mathbb{K} be a context such that $B(k)$ embeds into $\mathfrak{B}(\mathbb{K})$. Then $\mathbb{N}^c(k) \leq \mathbb{K}$.*

Proof. Let $(A_1, B_1), \dots, (A_k, B_k)$ be the atoms of $B(k)$ in $\mathfrak{B}(\mathbb{K})$. Similarly, denote its coatoms by $(C_1, D_1), \dots, (C_k, D_k)$ in such a way that $(A_i, B_i) \leq (C_j, D_j) \Leftrightarrow i \neq j$ for each i, j . Note that the sets A_i , as well as the sets D_i , are non-empty. Let $i \in [k]$. Since $(A_i, B_i) \not\leq (C_i, D_i)$, we may take an object/attribute pair $g_i \in A_i, m_i \in D_i$ with $g_i \not\leq m_i$. For every chosen object $g_i \in A_i$, one has that

$g_i I m_j$ for every $j \in [k]$ with $j \neq i$, because of $(A_i, B_i) \leq (C_j, D_j)$, which implies $B_i \supseteq D_j$. Consequently, k distinct objects g_i (as well as k distinct attributes m_i) were chosen. Combining both relations results in $g_i I m_j \Leftrightarrow i \neq j$ for each $i \in [k]$, that is, the objects and attributes g_i, m_i form a contranominal scale in \mathbb{K} . \square

Definition 1. Let (G, M, I) be a formal context. A set $S \subseteq G$ is said to be a **minimal generator** (of the extent S'') if $T'' \neq S''$ for every proper subset $T \subsetneq S$. The set of all minimal generators of a context \mathbb{K} will be denoted by $\text{MINGEN}(\mathbb{K})$.

Observation: In contexts with finitely many objects, every extent has at least one minimal generator. Clearly, two different extents cannot share one same minimal generator. Thus, the upper bound $|\mathfrak{B}(\mathbb{K})| \leq |\text{MINGEN}(\mathbb{K})|$ holds for contexts with finite object sets.

The problem of computing exactly the number of concepts does not admit a polynomial-time algorithm, unless $P=NP$. This was shown by Kuznetsov [5]: more precisely, this counting problem is $\#P$ -complete. However, there are results which establish upper bounds for the number of concepts: see for example [1–3, 6, 7].

3 The upper bound

Our investigations were inspired by a result of Prisner, who gave the first upper bound regarding contranominal-scale free contexts. The original version is in graph theoretic language. Reformulated it reads as follows:

Theorem 1 (Prisner [6]). Let $\mathbb{K} = (G, M, I)$ be a $\mathbb{N}^c(k)$ -free context. Then,

$$|\mathfrak{B}(\mathbb{K})| \leq (|G||M|)^{k-1} + 1.$$

In this section we will show an improvement of Theorem 1. For that, we will relate minimal generators with contranominal scales. The first step towards this is the equivalence shown in Proposition 2. Note that, since derivation operators are antitone, the \neq symbol may be substituted by \supsetneq .

Proposition 2. Let (G, M, I) be a formal context. A set $S \subseteq G$ is a minimal generator if and only if for every $g \in S$, it holds that $(S \setminus \{g\})' \neq S'$.

Proof. We will show the two equivalent contrapositions. If $(S \setminus \{g\})' = S'$, then, of course, $(S \setminus \{g\})'' = S''$, and S is not a minimal generator. For the converse, suppose that S is not a minimal generator, and take a proper subset T of S with $T'' = S''$. Note that $T'' = S''$ implies $T' = S'$. Let $g \in S \setminus T$. On one hand, $(S \setminus \{g\}) \subseteq S$ implies $(S \setminus \{g\})' \supseteq S'$. On the other hand, $(S \setminus \{g\}) \supseteq T$ implies $(S \setminus \{g\})' \subseteq T' = S'$. Combining both yields $(S \setminus \{g\})' = S'$. \square

The next proposition relates minimal generators and contranominal scales.

Lemma 1. *Let $\mathbb{K} = (G, M, I)$ be a context and $A \subseteq G$. There exists a contranominal scale $\mathbb{K}_1 \leq \mathbb{K}$ having A as its object set if and only if A is a minimal generator. In particular, if G is finite:*

$$\max_{A \subseteq G} \{|A| : A \text{ is a minimal generator}\} = \max\{k \in \mathbb{N} : \mathbb{N}^c(k) \leq \mathbb{K}\}.$$

Proof. Suppose that A is a minimal generator and let $g \in A$. By Proposition 2, one has that $(A \setminus \{g\})' \supseteq A'$. Hence, there exists an attribute m with gFm and hIm for every $h \in A \setminus \{g\}$. Clearly, two different objects $g_1, g_2 \in A$ cannot give rise to the same attribute m , since the two pairs of conditions g_iFm and hIm for every $h \in A \setminus \{g_i\}$ cannot be satisfied simultaneously ($i = 1, 2$). Thus, there exists an injection $\iota : A \rightarrow M$ with $gF\iota(g)$, $hI\iota(g)$ for each $g \in A$ and each $h \in A \setminus \{g\}$. By setting $N = \iota(A)$, one has that $(A, N, I \cap (A \times N))$ is a contranominal scale. For the converse, let $\mathbb{K}_1 = (S, S, \neq) \leq \mathbb{K}$ be a contranominal scale and let $g \in S$. Clearly, $g \notin S'$. Moreover, $g \in (S \setminus \{g\})'$. This amounts to $(S \setminus \{g\})' \supseteq S'$ for each $g \in S$. By Proposition 2, the set S is a minimal generator. \square

A consequence of Lemma 1 is the following, which is an improvement of the order of $k! \cdot |M|^k/k$ of Prisner's bound.

Theorem 2. *Let $\mathbb{K} = (G, M, I)$ be a $\mathbb{N}^c(k)$ -free context with finite G . Then:*

$$|\mathfrak{B}(\mathbb{K})| \leq |\text{MINGEN}(\mathbb{K})| \leq \sum_{i=0}^{k-1} \binom{|G|}{i}.$$

In particular, if $k \leq \frac{|G|}{2}$:

$$|\mathfrak{B}(\mathbb{K})| \leq k \cdot \frac{|G|^{k-1}}{(k-1)!}.$$

Proof. Lemma 1 guarantees that \mathbb{K} does not have any minimal generator of cardinality greater or equal to k . The sum above is the number of subsets of G having cardinality at most $k-1$. \square

Definition 2. *We denote by $f(n, k)$ the upper bound in Theorem 2:*

$$f(n, k) := \sum_{i=0}^{k-1} \binom{n}{i}.$$

The upper bound in Theorem 2 for $f(n, k)$ gets worse as k gets close to $\frac{|G|}{2}$. Tighter upper bounds for the sum of binomial coefficients may be found in [9].

4 Sharpness: preparatory results

The following property of $f(n, k)$ is needed for the next two sections.

Proposition 3. *The function $f(n, k)$ satisfies the following identity:*

$$f(n, k) = f(n - 1, k - 1) + f(n - 1, k).$$

Proof. This follows from a standard binomial identity: $f(n - 1, k) + f(n - 1, k - 1) = \sum_{i=0}^{k-1} \binom{n-1}{i} + \sum_{j=0}^{k-2} \binom{n-1}{j} = 1 + \sum_{i=1}^{k-1} \binom{n-1}{i-1} + \sum_{j=1}^{k-1} \binom{n-1}{j} = 1 + \sum_{i=1}^{k-1} \binom{n}{i} = f(n, k)$. \square

Consider a finite lattice L . It is well known that every element $x \in L$ is the supremum of some subset of $J(L)$: for example, $x = \bigvee J_x$. We call such a subset a *representation of x through join-irreducible elements* (for brevity, we may say a *representation through irreducibles of x* or even only a *representation of x*). A representation $S \subseteq J(L)$ of x is called *irredundant* if $\bigvee(S \setminus \{y\}) \neq x$ for every $y \in S$. Of course, every $x \in L$ has an irredundant representation, but it does not need to be unique. Note that irredundant representations are precisely minimal generators when one takes the standard context of L , $(J(L), M(L), \leq)$. Indeed, in that formal context, the closure of object sets corresponds to the supremum of join-irreducible elements of L . For an element $x \in L$, there may exist elements in J_x which belong to every representation of x : the so-called *extremal points*. An element $z \in J_x$ is an *extremal point of x* if there exists a lower neighbor y of x such that $J_y = J_x \setminus \{z\}$. Every representation of x must contain every extremal point z of x since, in this case, the supremum $\bigvee(J_x \setminus \{z\})$ is strictly smaller than x (and is actually covered by x).

In Section 5 we shall construct finite lattices for which every element has exactly one irredundant representation. It turns out that, in the finite case, these lattices are precisely the meet-distributive lattices. This is implied by Theorem 44 of [4], which actually gives information about the unique irredundant representation as well: a finite lattice L is meet-distributive if and only if for every $x \in L$ the set E_x of all extremal points of x is a representation of x (and E_x is, therefore, the unique irredundant representation of x , since every representation of x must contain E_x). Proposition 4 provides a characteristic property for the finite case which will be used in our constructions.

Proposition 4. *Let L be a finite lattice. The following assertions are equivalent:*

- i) L is meet-distributive.*
- ii) Every element $x \in L$ is the supremum of its extremal points.*
- iii) For every $x, y \in L$ with $x \prec y$, it holds that $|J_y \setminus J_x| = 1$.*

Proof. The equivalence between *i)* and *ii)* may be found in Theorem 44 of [4]. Let $x \in L$ and define $E_x = \{z \in J_x \mid z \text{ is an extremal point of } x\}$. We now show that *ii)* implies *iii)*. Let $y \in L$ with $y < x$. This implies $J_y \subsetneq J_x$. The set J_y does not contain E_x , because this would force $y \geq x$. Therefore, $y = \bigvee J_y$ is upper bounded by some element in the set $U = \{\bigvee(J_x \setminus \{z\}) \mid z \in E_x\}$ (note that $x \notin U$). Hence, every lower neighbor of x has a representation of the form $(J_x \setminus \{z\})$ with $z \in E_x$. Now we show that *iii)* implies *ii)*. Define $y = \bigvee E_x$ and suppose by contradiction that $y < x$. Then, there exists an element z such

that $y \leq z \prec x$ and $J_z \supseteq E_x$. But then, $z \prec x$ implies $J_x \setminus J_z = \{w\}$ for some $w \in J(L)$, which means that w is an extremal point of x . This contradicts the fact that E_x contains all extremal points of x . \square

The next lemma will be useful in Section 5, when we shall change the perspective from lattices to contexts.

Lemma 2. *Let L be a finite lattice. If L is $B(k)$ -free, then every element has a representation through join-irreducibles of size at most $k-1$. The converse holds if L is meet-distributive.*

Proof. Let $\mathbb{K} = (J(L), M(L), \leq)$ be the standard context of L . We identify the elements of L with the extents of \mathbb{K} via $x \mapsto J_x$. Suppose that L is $B(k)$ -free. Then, \mathbb{K} is $N^c(k)$ -free. Let A be an arbitrary extent of \mathbb{K} and S a minimal generator of A . Then, by Lemma 1, it follows that $|S| \leq k-1$. Since $A = S'' = \bigvee S$, we have the desired representation. Now, suppose that $|J_y \setminus J_x| = 1$ holds for every $x, y \in L$ with $x \prec y$ (cf. Proposition 4). To prove the converse, we suppose that $B(k)$ embeds into L and our goal is to show that some $x \in L$ does not have any representation with fewer than k elements of $J(L)$. Now, since $B(k)$ embeds into L , Proposition 1 implies that $N^c(k)$ is a subcontext of \mathbb{K} . Applying Lemma 1, we have that there exists a minimal generator $S \subseteq J(L)$ with $|S| = k$. Equivalently, S is an irredundant representation of the element S'' of L . By Proposition 4, S is the unique irredundant representation of S'' . Therefore, S'' cannot be expressed as the supremum of fewer than k join-irreducible elements. \square

5 Sharpness: construction of extremal lattices

In this section, we will consider only finite lattices. Our objective is to construct lattices which prove that the bound in Theorem 2 is sharp.

Definition 3. *For positive integers n and k , we call a lattice **(n, k) -extremal** if it has at most n join-irreducible elements, is $B(k)$ -free, and has exactly $f(n, k)$ elements.*

It is clear that every $(n, 1)$ -extremal lattice is trivial, i.e., the lattice with one element. To construct (n, k) -extremal lattices with larger k , we will use an operation which we call *doubling*.

Definition 4. *Let L be an ordered set and $K \subseteq L$. The **doubling of K in L** is defined to be $L[K] = L \cup \dot{K}$, where \dot{K} is a disjoint copy of K , i.e., $\dot{K} \cap L = \emptyset$. The order in $(L[K], \leq')$ is defined as follows:*

$$\leq' = \leq \cup \{(x, \dot{y}) \in L \times \dot{K} \mid x \leq y\} \cup \{(\dot{x}, \dot{y}) \in \dot{K} \times \dot{K} \mid x \leq y\}.$$

We will employ the notation \dot{x} to denote the image under doubling of an element $x \in K$. Note that $x \prec \dot{x}$ for every $x \in K$, and that \dot{x} is the only upper neighbor of x in \dot{K} . When L is a set family $\mathcal{C} \subseteq \mathcal{P}(G)$, then the diagram of $L[K]$ can be easily depicted: the doubling $\mathcal{C}[\mathcal{D}]$ (with $\mathcal{D} \subseteq \mathcal{C}$) corresponds to the set family $\mathcal{C} \cup \{D \cup \{g\} \mid D \in \mathcal{D}\}$, where $g \notin G$ is a new element. Figure 2 illustrates three doubling operations. The first one is the doubling of the chain $\{\emptyset, \{2\}, \{1, 2\}\}$ inside the closure system $\mathcal{C}_1 = \mathcal{P}(\{2\})$, resulting in \mathcal{C}_2 . The (*a fortiori*) closure systems \mathcal{C}_3 and \mathcal{C}_4 are obtained by doubling, respectively, the chains $\{\emptyset, \{3\}, \{2, 3\}, \{1, 2, 3\}\}$ and $\{\emptyset, \{2\}, \{2, 3\}, \{1, 2, 3\}\}$ inside \mathcal{C}_2 .

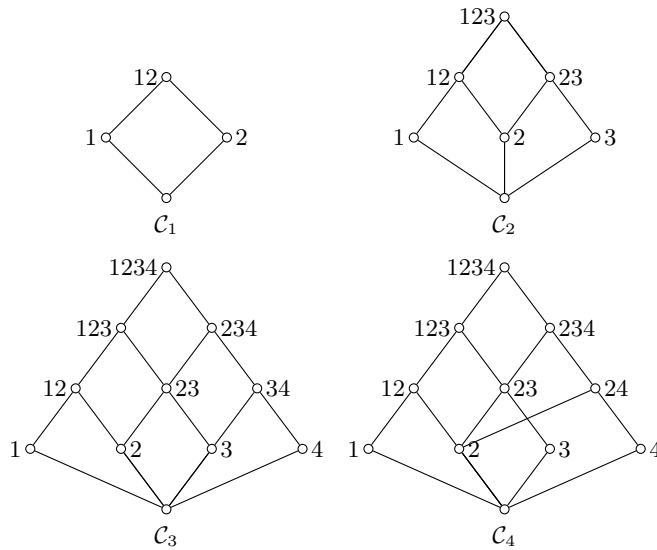


Fig. 2: Doubling chains inside closure systems

Since we are interested in constructing lattices, it is vital to guarantee that the doubling operation produces a lattice. By a *meet-subsemilattice* of a lattice L is meant a subset K of L , endowed with the inherited order, such that $x \wedge y \in K$ holds for every $x, y \in K$. It is called *topped* if $1_L \in K$.

Proposition 5. *If K is a topped meet-subsemilattice of a lattice L , then $L[K]$ is a lattice.*

Proof. Let $x, y \in L[K]$. If both x and y belong to L , then clearly $x \wedge y$ and $x \vee y$ belong to $L \subseteq L[K]$. Suppose that only one among x and y , say x , belongs to L . Then $y = \dot{z}$ with $z \in K$. We have that $x \wedge y = x \wedge z \in L \subseteq L[K]$ because of $x \not\leq 0_K$ and $y = z \vee 0_K$. For the supremum, set $S = \{w \in K \mid w \geq x, w \geq z\}$ and $u = \bigwedge S$. Note that the fact that K is topped causes $S \neq \emptyset$. Since K is a meet-subsemilattice, we have that $u \in K$. It is clear that u is the least upper

bound of x and z which belongs to K . Therefore, \dot{u} is the least upper bound of x and y , because of $0_K \not\leq u$ and $y = z \vee 0_K$. The remaining case is $x, y \in \dot{K}$ for which, clearly $x \wedge y$ exists. Moreover, writing $x = \dot{t}, y = \dot{z}$ with $t, z \in K$ and setting $S = \{w \in K \mid w \geq t, w \geq z\}$ as well as $u = \bigwedge S$ make clear that $\dot{u} = x \vee y$. \square

When extrinsically considered, topped meet-subsemilattices are lattices. This is compatible with the proof of Proposition 5, where the supremum and infimum of two elements in \dot{K} may be easily verified to belong to \dot{K} : that is, \dot{K} is actually a sublattice of $L[K]$.

A suborder K of an ordered set L is called *cover-preserving* if $x \prec_K y$ implies $x \prec_L y$ for every $x, y \in K$. This property plays a key role by preserving meet-distributivity under a doubling operation:

Proposition 6. *Let L be a meet-distributive lattice and let K be a cover-preserving, topped meet-subsemilattice of L . Then, $L[K]$ is a meet-distributive lattice.*

Proof. The fact that $L[K]$ is a lattice comes from Proposition 5. Every element $\dot{x} \in \dot{K}$ has one lower neighbor in K , namely, x . Thus, the total number of lower neighbors of \dot{x} is one only if x does not cover any element in K , that is, $x = 0_K$. Therefore, 0_K is the only join-irreducible of $L[K]$ which is not a join-irreducible of L . Let $x, y \in L[K]$ with $x \prec_{L[K]} y$. We use $J'_{(\cdot)}$ to denote our J -notation in $L[K]$ and $J_{(\cdot)}$ in L . If $x, y \in L$, then clearly $J'_y = J_y$ and $J'_x = J_x$, which results in $|J'_y \setminus J'_x| = |J_y \setminus J_x| = 1$. If $x, y \notin L$, then $x = \dot{z}$ and $y = \dot{w}$ with $z, w \in K$ and $z \prec_K w$. From the fact that K is cover-preserving, we conclude that $z \prec_L w$. Because L is meet-distributive, it follows that $|J_w \setminus J_z| = 1$. Clearly one has $J'_x = J_z \cup \{0_K\}$ and $J'_y = J_w \cup \{0_K\}$, which yields $|J'_y \setminus J'_x| = 1$. For the remaining case, one has necessarily $x \in L$ and $y \notin L$. In these conditions, $x \prec y$ results in $y = \dot{x}$ and, therefore, $J'_y = J'_x \cup \{0_K\}$, implying $|J'_y \setminus J'_x| = 1$. \square

Proposition 7 is the first assertion about extremal meet-subsemilattices. We note that the set of join-irreducible elements of a meet-subsemilattice K of a lattice L is not the same as the set $J(L) \cap K$. Therefore, what is meant by an (n, k) -*extremal meet-subsemilattice* of L is precisely the following: a lattice K which is (n, k) -extremal and a meet-subsemilattice of L as well.

Observe that chains with $n + 1$ elements are precisely the $(n, 2)$ -extremal lattices. Proposition 7 illustrates, in particular, that an $n + 1$ element chain may be seen as the result of a doubling operation on an n element chain, provided that the doubling operation is performed with respect to the trivial topped meet-subsemilattice $\uparrow 1$, which is $(n, 1)$ -extremal.

Proposition 7. *Let L be an $(n - 1, k)$ -extremal lattice with $n, k \geq 2$. Suppose that K is a topped, $(n - 1, k - 1)$ -extremal meet-subsemilattice. If $L[K]$ is $B(k)$ -free, then it is an (n, k) -extremal lattice.*

Proof. Proposition 5 guarantees that $L[K]$ is indeed a lattice. As in the proof of Proposition 6, we have that $J(L[K]) = J(L) \cup \{0_K\}$ and in particular, $L[K]$ has at most n join-irreducible elements. The claim that $L[K]$ has $f(n, k)$ elements follows from Proposition 3. \square

It is clear now how $(n, 2)$ -extremal lattices can be obtained by doubling a trivial meet-subsemilattice of an $(n - 1, 2)$ -extremal lattice. The succeeding propositions and lemmas aim particularly towards a generalization of this operation: the doubling of topped, $(n - 1, k - 1)$ -extremal meet-subsemilattices inside $(n - 1, k)$ -extremal lattices, yielding (n, k) -extremal lattices for $k \geq 3$.

Proposition 8. *Suppose that L is an (n, k) -extremal lattice. Then, for every $S, T \subseteq J(L)$ with $|S|, |T| \leq k - 1$:*

$$\bigvee S = \bigvee T \Rightarrow S = T.$$

Moreover, if $k \geq 2$ then $|J(L)| = n$.

Proof. We may suppose $k \geq 2$ since the assertion holds trivially for $k = 1$. Lemma 2 guarantees that every element x of L has a representation of size at most $k - 1$. Therefore $L = \{\bigvee S \mid S \subseteq J(L), |S| \leq k - 1\}$. Because of $k \geq 2$ and the fact that $|L| = f(n, k)$ is also the number of subsets of $[n]$ having at most $k - 1$ elements, one has $|J(L)| \geq n$. In fact equality must hold, because L has at most n join-irreducible elements. As a consequence of $|J(L)| = n$ and $|L| = f(n, k)$, we have that no two sets $S, T \subseteq J(L)$ with $S \neq T$ may lead to the same supremum $\bigvee S = \bigvee T$. \square

Chains are the only extremal lattices which are not atomistic, as a consequence of the next lemma.

Lemma 3. *Suppose that L is an (n, k) -extremal lattice with $k \geq 3$. Then L is atomistic and meet-distributive. In particular, the length of L equals the number of its atoms and there exists an atom which is an extremal point of 1_L .*

Proof. If L were not atomistic, there would exist two comparable join-irreducible elements, say, x, y with $x < y$. But then $x \vee y = y$, which contradicts Proposition 8. Suppose that L is not meet-distributive and take $x, y \in L$ with $x \prec y$ such that $A_y \setminus A_x$ has at least two elements. Clearly $x \neq 0_L$ and, therefore, $A_x \neq \emptyset$. Let $u, v \in A_y \setminus A_x$ be any two distinct elements. From $u \notin A_x$ follows that $x < x \vee u \leq y$ which, in turn, implies $x \vee u = y$. Similarly, $v \notin A_x$ implies $x < x \vee v \leq y$ which, in turn, implies $x \vee v = y$. Let $a \in A_x$. Now, $a \leq x$ and $x \parallel u$ imply $a \vee u = x \vee u = y$, as well as $a \leq x$ and $x \parallel v$ imply $a \vee v = x \vee v = y$. We obtain $a \vee u = a \vee v$, contradicting Proposition 8. Choosing a maximal chain $x_0 \prec x_1 \prec \dots \prec x_l$ in L and noticing that the sizes of the sets A_{x_i} grow by exactly one element make the two remaining claims clear. \square

Lemma 4 shows that non-trivial, extremal meet-subsemilattices are always cover-preserving and topped. These two properties will be useful to assure that a doubling $L[K]$ is a meet-distributive lattice.

Lemma 4. *Let L be an (n, k) -extremal lattice with $k \geq 3$ and suppose that K is an $(n, k - 1)$ -extremal subsemilattice of L . Then, K is cover-preserving and topped. If $k \geq 4$, then K and L are atomistic with $A(K) = A(L)$.*

Proof. In case that $k = 3$ then K is B(2)-free, that is, K is a chain. By Lemma 3, K must be a maximal chain in order to have $\sum_{i=0}^1 \binom{n}{i} = n + 1$ elements. Hence, $1_K = 1_L$. The maximality of K guarantees that K is cover-preserving. Now, suppose that $k \geq 4$. Again by Lemma 3, we have that both K and L are atomistic. Since K has n atoms, 0_K must be covered by n elements in L . But this is possible only if $0_L = 0_K$, because L also has n atoms. This forces $A(K) = A(L)$ as well as $1_L \in K$, because 1_L is the only element that upper bounds each $a \in A(K)$. To prove that K is cover-preserving, we apply Lemma 3 twice, obtaining $|A_y \setminus A_x| = 1$ for every $x, y \in K$ with $x \prec_K y$ as well as $|A_y \setminus A_x| = 1$ for every $x, y \in L$ with $x \prec_L y$. Both conditions hold simultaneously only if the implication $x \prec_K y \Rightarrow x \prec_L y$ holds, i.e., if K is cover-preserving. \square

A *complete meet-embedding* is a meet-embedding which preserves arbitrary meets, including $\bigwedge \emptyset$. As a consequence, the greatest element of one lattice gets mapped to the greatest element of the other. Images of complete meet-embeddings are topped meet-subsemilattices. This notion is required for the following simple fact, which aids us in the construction of sequences of (n, k) -extremal lattices with fixed n and growing k . In Proposition 9, the symbol $K[J]$ (for instance) means actually the doubling of the image of J under the corresponding embedding.

Proposition 9. *Suppose that J, K and L are lattices with complete meet-embeddings $\mathcal{E}_1 : J \rightarrow K$ and $\mathcal{E}_2 : K \rightarrow L$. Then, there exists a complete meet-embedding from $K[J]$ into $L[K]$.*

Proof. The fact that $K[J]$ and $L[K]$ are lattices comes from Proposition 5. Of course, there is an induced embedding from \dot{J} into \dot{K} , but for which we will use the same symbol \mathcal{E}_1 . The mapping $\mathcal{E}_3 : K[J] \rightarrow L[K]$ defined by $\mathcal{E}_3(x) = \mathcal{E}_1(x)$ for $x \in \dot{J}$ and $\mathcal{E}_3(x) = \mathcal{E}_2(x)$ for $x \in K$ may be checked as being a complete meet-embedding. \square

As mentioned after Proposition 7, we will make use of an operation which doubles an extremal meet-subsemilattice of an extremal lattice. The next theorem shows that the lattice produced by this operation is indeed extremal.

Theorem 3. *Let L be an $(n - 1, k)$ -extremal lattice with $n \geq 2$ and $k \geq 3$ and suppose that K is an $(n - 1, k - 1)$ -extremal meet-subsemilattice of L . Then, $L[K]$ is an (n, k) -extremal lattice.*

Proof. Lemma 3 guarantees that L is atomistic and meet-distributive. Moreover, Lemma 4 guarantees that K is cover-preserving and topped, so that, in particular, $L[K]$ is a meet-distributive lattice, as a consequence of Proposition 6. To prove that $L[K]$ is (n, k) -extremal it is sufficient to show that $L[K]$ is B(k)-free,

because of Proposition 7. We will do so by proving that every element of $L[K]$ has a representation through join-irreducibles of size at most $k - 1$. This indeed suffices because $L[K]$ is meet-distributive, so that Lemma 2 may be applied. Observe that $J(L[K]) = A(L) \cup 0_K$, since L is atomistic.

Suppose that $k = 3$. In this case K is a chain, and a maximal one because of Lemma 4. Let $x \in L[K]$. If $x \in L$, then Lemma 2 implies that $x = \bigvee S$ for some $S \subseteq A(L)$, $|S| \leq 2$ and the same representation may be used in $L[K]$. If $x \notin L$, then $x = \dot{y} = y \vee 0_K$ for some $y \in K$. If $y \in A(L)$, we are done. Otherwise, take $z, w \in A(L)$ such that $z \vee w = y$ and thus $x = z \vee w \vee 0_K$. Exactly one among z and w belong to K . Without loss of generality, let it be z . Then, it is clear that $0_K \prec z \vee 0_K$ and that $z \vee 0_K \prec w \vee 0_K$, since there exists only one element covering 0_K . Hence, $x = z \vee w \vee 0_K = w \vee 0_K$.

Suppose that $k \geq 4$. As noted after Proposition 5 one has that K is, by itself, a lattice. Moreover, Lemma 4 guarantees that K is atomistic with $A(L) = A(K)$. Let $x \in L[K]$. If $x \in L$ then, in L , $x = \bigvee S$ for some $S \subseteq A(L) \subseteq J(L[K])$ with $|S| \leq k - 1$, because of Lemma 2 and the fact that L is $B(k)$ -free. Of course, S is also a representation of x in $L[K]$. If $x \notin L$, then $x = \dot{y}$ for some $\dot{y} \in \dot{K}$. Since K is $B(k - 1)$ -free, it follows that, in K , $y = \bigvee S$ for some $S \subseteq A(K) \subseteq J(L[K])$ with $|S| \leq k - 2$, once again as a consequence of Lemma 2. Clearly, in $L[K]$, one has $\dot{y} = \bigvee \dot{S} = 0_K \vee \bigvee S$, where the last equality follows from the fact that $\dot{z} = z \vee 0_K$ for every $z \in S$. Thus, we have a representation of $\dot{y} = x$ through no more than $k - 1$ join-irreducible elements of $L[K]$. \square

Corollary 1 describes how (n, k) -extremal lattices can be non-deterministically constructed. In particular, the upper bound present in Theorem 2 is the best possible.

Corollary 1. *For every n and k , there exists at least one (n, k) -extremal lattice.*

Proof. Define a partial function Φ satisfying

$$\Phi : \mathbb{N}^* \times \mathbb{N}^* \rightarrow \mathcal{L}$$

$$(n, k) \mapsto \begin{cases} ([n], \subseteq), & \text{if } k = 1. \\ (\{\emptyset, \{1\}\}, \subseteq), & \text{if } k \geq 2, n = 1. \\ \Phi(n - 1, k)[\mathcal{E}(\Phi(n - 1, k - 1))], & \text{if } n, k \geq 2 \text{ and there exists a} \\ & \text{complete meet-embedding} \\ & \mathcal{E} : \Phi(n - 1, k - 1) \rightarrow \Phi(n - 1, k). \end{cases}$$

where \mathcal{L} is the class of all lattices. We prove by induction on n that $\Phi(n, k)$ is a total function. The cases $n = 1$ and $n = 2$ are trivial. Let $n \in \mathbb{N}$ with $n \geq 3$ and suppose that $\Phi(n - 1, k)$ is defined for every $k \in \mathbb{N}^*$. Let $k \in \mathbb{N}$, $k \geq 2$. By the induction hypothesis, the values $\Phi(n - 1, k)$ and $\Phi(n - 1, k - 1)$ are defined. If $k = 2$, then $\Phi(n - 1, k - 1)$ is a trivial lattice and the existence of a complete meet-embedding into $\Phi(n - 1, k)$ is clear and, thereby, $\Phi(n, k)$ is defined. We therefore assume $k \geq 3$. By the definition of Φ , one has that $\Phi(n - 1, k) = \Phi(n -$

$2, k)[\mathcal{E}(\Phi(n-2, k-1))]$ and that $\Phi(n-1, k-1) = \Phi(n-2, k-1)[\mathcal{F}(\Phi(n-2, k-2))]$ for some pair of complete meet-embeddings \mathcal{E} and \mathcal{F} . Applying Proposition 9 with $\Phi(n-2, k-2)$, $\Phi(n-2, k-1)$ and $\Phi(n-2, k)$ results in the existence of a complete meet-embedding $\mathcal{G} : \Phi(n-1, k-1) \rightarrow \Phi(n-1, k)$, which yields that $\Phi(n, k)$ is defined. Since k is arbitrary, every $\Phi(n, k)$ is defined. The (n, k) -extremality of each lattice can be proved by induction on n as well and by invoking Theorem 3. \square

Figure 3 depicts the diagrams of nine (n, k) -extremal lattices which are constructible by Corollary 1. It is true that, in general, (n, k) -extremal lattices are not unique up to isomorphism: note that the $(3, 3)$ and $(4, 3)$ -extremal lattices in Figure 3 are also present in Figure 2 as the lattices \mathcal{C}_2 and \mathcal{C}_3 . The lattice \mathcal{C}_4 , depicted in that same figure, is a $(4, 3)$ -extremal lattice which is not isomorphic to \mathcal{C}_3 . We shall, however, show in the next section that every extremal lattice arises from the construction described in Corollary 1.

6 Characterization of extremal lattices

In the last section, we constructed lattices whose sizes are exactly the upper bound present in Theorem 2. In this section, we will show that every lattice meeting those requirements must be obtained from our construction.

Lemma 5. *Let L be an atomistic lattice, a an atom and c a coatom with $A_c = A(L) \setminus \{a\}$. Then, the mapping $x \xrightarrow{\mathcal{E}} c \wedge x$ is a complete meet-embedding of $\uparrow a$ into $\downarrow c$ such that $\mathcal{E}(x) \prec x$ for every $x \in \uparrow a$.*

Proof. The fact that \mathcal{E} preserves non-empty meets is clear, since c is a fixed element. Also, 1_L is mapped to $c = 1_{\downarrow c}$, so that \mathcal{E} preserves arbitrary meets. Note that

$$A_{\mathcal{E}(x)} = A_{c \wedge x} = A_x \cap A_c = A_x \setminus \{a\}.$$

Hence, $\mathcal{E}(x) \prec x$ as well as $\mathcal{E}(x) \vee a = x$. The latter implies injectivity. \square

The next theorem shows that every extremal lattice is constructible by the process described in Corollary 1, and can be seen as a converse of that result.

Theorem 4. *Let L be an (n, k) -extremal lattice with $k \geq 3$. Then, $L = J \dot{\cup} K$ where J is an $(n-1, k)$ -extremal lattice and K is an $(n-1, k-1)$ -extremal lattice. Moreover, there exists a complete meet-embedding $\mathcal{E} : K \rightarrow J$ such that $\mathcal{E}(x) \prec x$ for every $x \in K$. In particular, $L \cong J[\mathcal{E}(K)]$.*

Proof. From Lemma 3, one has that L is atomistic and we may take an atom a which is an extremal point of 1_L , that is, $A(L) \setminus \{a\} = A_c$ with c being a coatom of L . Consider the lattices $J = \downarrow c$ and $K = \uparrow a$. Observe that $L = J \dot{\cup} K$ and let $\mathcal{E} : K \rightarrow J$ be a complete meet-embedding provided by Lemma 5. Clearly, J has $n-1$ atoms and is $B(k)$ -free, therefore, $|J| \leq f(n-1, k)$. Moreover, K must be $B(k-1)$ -free: indeed, if there existed $B \cong B(k-1)$ inside K , then

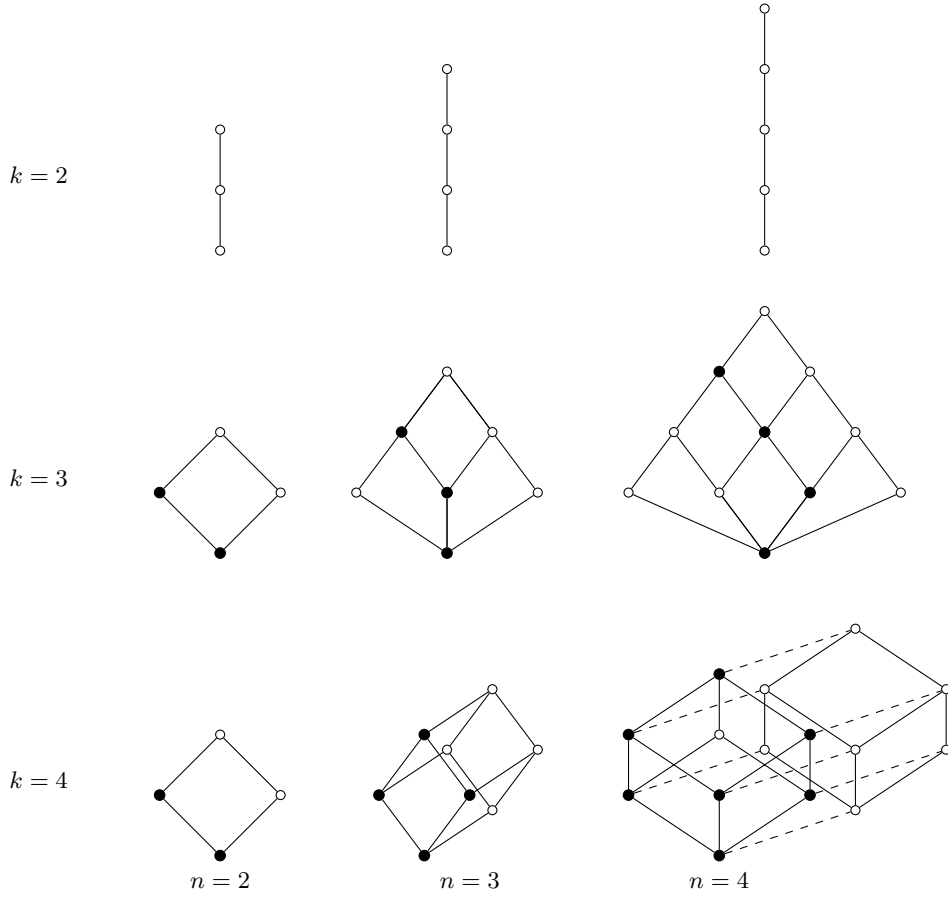


Fig. 3: Diagrams of (n, k) -extremal lattices with $2 \leq n, k \leq 4$. Elements shaded in black represent the doubled $(n - 1, k - 1)$ -extremal lattice.

$B \cup \mathcal{E}(B)$ would be a boolean lattice with k atoms inside J , which is impossible. The lattice K has at most $n - 1$ atoms, and consequently $|K| \leq f(n - 1, k - 1)$, since the function $n \mapsto \sum_{i=0}^{k-1} \binom{n}{i}$ is monotonic increasing. Now, we have that $|J| + |K| = |L| = f(n, k) = f(n - 1, k) + f(n - 1, k - 1)$, where the last equality follows from Proposition 3. Since $|J| \leq f(n - 1, k)$ and $|K| \leq f(n - 1, k - 1)$, those two inequalities must hold with equality. Therefore, J and K are, respectively, $(n - 1, k)$ and $(n - 1, k - 1)$ -extremal. \square

7 Conclusion and related work

We showed an upper bound for the number of minimal generators of a context which is sharp. Extremal lattices were constructed and also characterized. The rôle played by contranominal scales in formal contexts may be seen as analogous

as that of cliques in simple graphs, when one considers extremality with respect to the number of edges. The Sauer-Shelah Lemma [8] provides an upper bound which is similar to that of Theorem 2. This is not a coincidence, because it can be shown, not without some effort, that the condition of a concept lattice being $B(k)$ -free is equivalent to the family of its extents not shattering a set of size k . As for the sharpness of the bound (which we prove in Section 5), in our case it is non-trivial, whereas the sharpness for the result of Sauer and Shelah is immediate.

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